Document No:WG21 N4571Date:2015-05-08References:ISO/IEC PDTS 19841, SC22/WG21 N4488Reply To:Barry Hedquist <beh@peren.com><br/>INCITS/PL22.16 International Representative

#### **Record of Response: National Body Comments**

#### **ISO/IEC PDTS 19841**

#### **Technical Specification:** C++ Extensions for Transactional Memory

Attached is WG21 N4571, Record of Response to National Body Comments for ISO/IEC PDTS 19841, Technical Specification – C++ Extensions for Transactional Memory. Also attached is WG21 N4488, containing further details for the responses.

Document numbers referenced in this Record of Response are WG21 documents unless otherwise stated.

10ate:2015-05-06	Document: SC22/N 5019 WG21/N4517	Project: PDTS 19841
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MB/ NC <sup>1</sup>	Line number	Clause/ Subclause	Paragraph/ Figure/Table	Type of comment <sup>2</sup>	Comments	Proposed change	Observations of the secretariat
JP 1				pe la ac sc tra It re	Ve are concerned that there could be a erformance degradation even in an environment acking of transactional memory feature by dopting this technical specification, e.g. making ome standard libraries like <math.h> ransactional safe. : is the reason for our disapproval. If we can easonably confirm that there's no degradation, we will change our position to approval.</math.h>	Please make us sure there's no degradation.	REJECT See WG21 N4488 (Attch)
US 1				pr tra ef op pe	Memory ordering requirements of transactions are roblematically strict. Even empty or purely local ransactions have observable synchronization ffects and can usually not be removed by an ptimizing compiler. This introduces a erformance penalty when transactional library ode is reused in a clearly thread-local context.	Consider weakening ordering requirements to allow such optimizations.	ACCEPT See WG21 N4488 (Attch)
CA 1	N/A	N/A	N/A	ht	Request to add a Feature Test Macro _cpp_transactional_memory based on ttp://isocpp.org/std/standing-documents/sd-6- g10-feature-test-recommendations	The value of the macro will be the year and month of the release of the TS. It does not need any experimental or TS tag.	ACCEPT See WG21 N4488 (Attch)
CA 3	N/A	4.3 [conv.func]	Para 1	H tra te	<pre>Make helper functions in 20.2 transaction-safe. Here is an example where std::move is not ransaction-safe emplate <class t=""> void safe_swap(T &amp;a, T &amp;b) transaction_safe { atomic_commit { using std::move; T temp = move(a); // Note that std::move is ot transaction-safe according to draft, but it hould be a = move(b); b = move(temp);</class></pre>	Add std::move and other utilities in 20.2 to be transaction_safe.	ACCEPT See WG21 N4488 (Attch)

MB = Member body / NC = National Committee (enter the ISO 3166 two-letter country code, e.g. CN for China; comments from the ISO/CS editing unit are identified by \*\*)
Type of comment: ge = general te = technical ed = editorial

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				un } int { i	<pre>} } mplate <class t=""> void apply(T &amp;a, T &amp;b, void f(T&amp;,&amp;T)) {     f(a,b); // Ok     assert(f == safe_swap<int>); // result     nspecified according to 5.10, paragraph 2, right? } t main() int x = 2, y = 3; apply(x, y, safe_swap<int>); // Ok even though ansaction_safe is lost</int></int></class></pre>		
CA 2	N/A	5.2.2 [expr.call]	Para 1	A d ev [st in de ov It i fur tx_ dy ob	his addition states: call to a virtual function that is valuated within a synchronized (6.9 tmt.sync]) or atomic block (6.10 [stmt.tx]) results undefined behavior if the virtual function is eclared transaction_safe_noinherit and the final verrider is not declared transaction_safe. is Undefined Behavior if you call into a virtual inction declared as tx_safe_noinherit but it is not _safe in the final overrider. This ensures that the ynamic call is safe, no matter what the dynamic oject is since tx_safe_noinherit gives no uch guarantee.	Please fix for synchronized block so that it is not part of this requirement. Suggested wording: A call to a virtual function that is evaluated within a synchronized (6.9 [stmt.sync]) or an atomic block (6.10 [stmt.tx]) results in undefined behavior if the virtual function is declared transaction_safe_noinherit and the final overrider is not declared transaction_safe.	ACCEPT See WG21 N4488 (Attch)

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					synchronized block because these can call x_unsafe functions.		
FI 1		8	4	tı v ti n tı	t seems that the applicability of ransaction_safe_noinherit is likely going to be vider in the future than just in virtual functions. If hat wider applicability appears, new keywords need to be added. Generalizing the name ransaction_safe_noinherit would possibly avoid hat problem.	Rename transaction_safe_noinherit to transaction_safe_dynamic. Transaction safety of calls to such functions is ultimately a runtime property, hence _dynamic seems like a suitable suffix.	ACCEPT See WG21 N4488 (Attch)
JP 2		8.4.4	1	b	A function-local static variable initialization should be transactional-unsafe. The initialization in an atomic execution needs to be synchronized with non-atomic executions.	Add "a function-local static variable initialization" in the list of conditions for a transactional-unsafe statement .	REJECT See WG21 N4488 (Attch)
CA 4	n/a	8.4.4 [dcl.fct.def.t x]	After Para 1, bullet 5	tti s tti « ;; c c c s li s {	n the first sequence of dash bullets ( ) indicating ransaction-unsafe expressions, the fifth one states «an implicit call of a non-virtual function hat is not transaction_safe». I wonder why the kimplicit» call is being explicitly (sorry for the pun! ) ) specified, as it seems to me that an explicit call to a non-virtual function would yield the same consequences. Unless I'm missing out on something, an implicit call could be something ike: struct B int f(); // not transaction_safe, not virtual virtual ~B() = default; ;	This seems a possible confusion for other user, please clarify.	REJECT See WG21 N4488 (Attch)

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				ii { };	nt g() return f() + // implicit call? this->f() + // explicit call? B::f(); // explicit call?		

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ISO/IEC JTC 1/SC 22/WG 21 N4488 revises N4410 Jens Maurer 2015-05-07

## N4488: Responses to PDTS comments on Transactional Memory, version 2

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with other members of the transactional memory study group (SG5), including (in alphabetical order): Hans Boehm, hboehm@google.com

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## Introduction

This paper presents the proposed responses to N4396 "National Body Comments, ISO/IEC PDTS 19841, C++ Extensions for Transactional Memory".

## **Changes compared to N4410**

- fix a typo in the spelling of the feature test macro (two underscores in the middle)
- remove "L" suffix in value of feature test macro
- declval does not need to be transaction-safe, since it is never odr-used
- rework section CA 1

## JP 1: Performance degradation

#### REJECT

The specification of transaction safety ensures that it is possible to compile code so that whether a given function call is executed in transaction context or outside of transaction context can be determined at compile time. Therefore, existing code that is executed outside of transactions and that does not use any of the transactional memory constructs will execute as before. The performance of code that actually uses transactions will depend on the available hardware support, similar to the fact that the performance of mutexes vs. accesses to atomic variables depends on a number of hardware and other factors.

A conservative approach was chosen for mandating the transaction-safety of standard library functions. Functions that conceivably access global state are not touched. In particular, the functions in the header <math.h> were intentionally not made transaction-safe in this Technical Specification, because the interaction of transactional memory with accesses to a potentially global rounding mode setting was deemed to require further study. As an exception, based on early user feedback, memory allocation is mandated to be transaction-safe, although it might access the global free store. Implementation experience shows that this does not negatively impact the performance of non-transactional executions.

# **US 1: Relax synchronization to allow optimizations on local transactions**

#### ACCEPT

In section 1.10 intro.multithread, change the added paragraph 9 as follows:

There is a global total order of execution for all outer blocks. If, in that total order, T1 is ordered before T2,

- no evaluation in T2 happens before any evaluation in T1 and
- **if T1 and T2 perform conflicting expression evaluations,** then the end of T1 synchronizes with the start of T2.

## CA 1: Feature Test Macro

#### ACCEPT

Change in section 1.3 [general.references]:

... Beginning with section 1410 below, all clause and section numbers, titles, and symbolic references in [brackets] refer to the corresponding elements of the C++ Standard. Sections 1.1 through 150 of this Technical Specification are introductory material and are unrelated to the similarly-numbered sections of the C++ Standard.

Change in section 1.4 [intro.compliance]:

Conformance requirements for this specification are the same as those defined in <u>section</u> 1.4 [intro.compliance] <u>of the C++ Standard</u>. [Note: Conformance is defined in terms of the behavior of programs. -- end note ] Add a new section 1.5 [intro.features]:

#### 1.5 Feature testing [intro.features]

An implementation that provides support for this Technical Specification shall define the feature test macro in Table 1.

#### Table 1 -- Feature Test Macro

Name	Value	Header
cpp_transactional_memory	201505	predeclared

## CA 3: Make helper functions transaction-safe

ACCEPT

Add the following to the appropriate sections:

Add in 20.2 [utility] after the synopsis: <u>A function in this section is transaction-safe if all required operations are</u> <u>transaction-safe.</u> In 20.2.4 [forward], add "transaction\_safe" to the declaration of all functions.

## CA 2: Virtual function calls in synchronized blocks

#### ACCEPT

Change the added text in 5.2.2 [expr.call] paragraph 1:

A call to a virtual function that is evaluated within a synchronized (6.9 [stmt.sync]) or an atomic block (6.10 [stmt.tx]) results in undefined behavior if the virtual function is declared transaction\_safe\_noinherit transaction\_safe\_dynamic and the final overrider is not declared transaction\_safe.

#### $FI \ 1: Rename \ \texttt{transaction\_safe\_noinherit} \ to \ \texttt{transaction\_safe\_dynamic}$

#### ACCEPT

Change all mentions of transaction\_safe\_noinherit to transaction\_safe\_dynamic, including sections 2.11 [lex.name], 5.2.2 [expr.call] (see also CA 2, above), clause 8 [dcl.decl], 8.3.5 [dcl.fct], 10.3 [class.virtual], 18.6.2.1 [bad.alloc], 18.6.2.2 [new.badlength], 18.7.2 [bad.cast], 18.7.3 [bad.typeid], 18.8.1 [exception], 18.8.2 [bad.exception], and 19.2 [std.exceptions].

## JP 2: Initialization of function-local static variables

REJECT

We agree that initialization of function-local statics should be atomic with respect to both transactional and non-transactional uses. We do not believe that the specification as drafted, taking into consideration the requirements on transaction safety, necessitates any additional overhead on the non-transactional code path once the initialization is complete.

## CA 4: Redundant case for transaction-unsafe expressions

#### REJECT

The specification is carefully crafted to ensure that calls through function pointers or member function pointers fall into the sixth bullet. Omitting "implicit" in the fifth bullet would (arguably) defeat that purpose.